

NINETY-SIXTH INTERNATIONAL ATLANTIC ECONOMIC CONFERENCE

PRESIDENTIAL ADDRESS

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PRESIDENT

INTERNATIONAL ATLANTIC ECONOMIC SOCIETY

“On Digital Currencies”

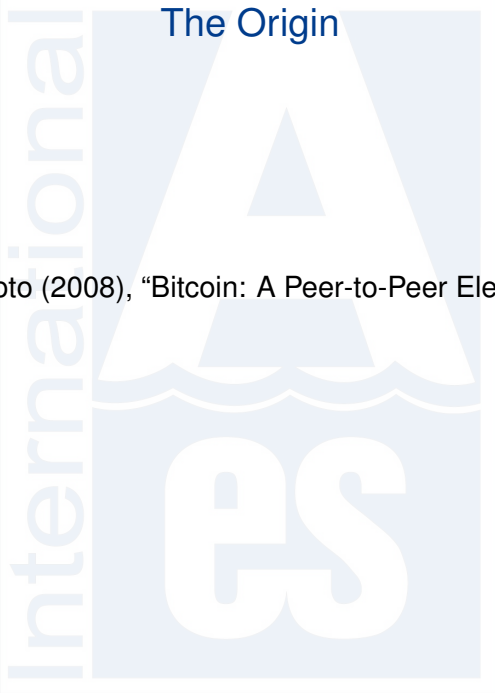
5-8 OCTOBER 2023

This talk:

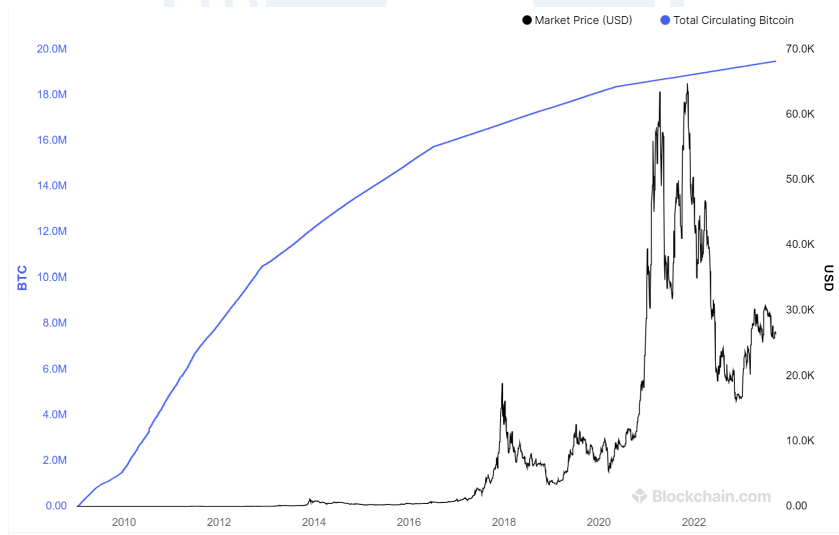
- 1 Theory of Cryptocurrency Pricing: **Schilling-Uhlig**, “**Some Simple Bitcoin Economics**”, JME 2019.
- 2 Central bank digital currencies: **Schilling - Fernández-Villaverde - Uhlig**, “**CBDC: when Price and Bank Stability Collide**”, WP.

The Origin

Satoshi Nakamoto (2008), "Bitcoin: A Peer-to-Peer Electronic Cash System."



Bitcoin Quantity and Price, 2009-01-09 to 2023-09-26



Source: <https://www.blockchain.com/charts/total-bitcoins>

Schilling - Uhlig, “Some Simple Bitcoin Economics”

Key Questions:

- 1 What determines the Bitcoin price? $P > NPV(\text{Dividends}) = 0$?
- 2 Can Bitcoin serve as medium of exchange, despite price volatility?
- 3 What are monetary policy implications?

Key Insights:

- 1 A **novel model** of an endowment economy with two intrinsically worthless currencies (Dollar, Bitcoin) as medium of exchange.
- 2 “Fundamental pricing equation”. Special case: **Bitcoin price is martingale**.
 - ▶ KAREKEN AND WALLACE (1981)
 - ▶ MANUELLI AND PECK (1990)
- 3 **“No speculation” theorem**.
- 4 **Volatility does not invalidate medium-of-exchange function**.
- 5 **Monetary policy implications:**
 - ▶ **Bitcoin block rewards are not a tax on Bitcoin holders: they are financed with a Dollar tax.**

The Schilling - Uhlig (2019) model

- Time: discrete, infinite $t = 0, 1, 2, \dots$
- Randomness: θ_t , at beginning of period. History: θ^t .
- One perishable consumption good per period.
- Two monies: Bitcoins B_t and Dollars D_t (aggregates).
- A central bank steers quantity of Dollars per lump sum transfers:
 $D_{t+1} = D_t + \tau_{t+1}$, $\tau_{t+1} \in \mathbb{R}$. Goal: exogenous price path P_t .
- Bitcoin quantity: deterministic
 $B_{t+1} = B_t + A_t$, $A_t \geq 0$ (endowment or “mining”)
- P_t price of consumption good in \$: exogenous.
- $Q_t = Q(\theta^t)$ price of Bitcoins in \$: **endogenous**.
- Two types of agents: “red” and “green”. “Red” agents consume in odd periods and have endowments in even periods. “Green” agents other way around.
- Goods are traded for monies. Agents do not need to spend all money (“hodlers”) or accept all money. **But**: “No speculation” result: they will!

The Fundamental Pricing Equation

Compare to Kareken-Wallace (1981), Manuelli-Peck (1990)

Proposition 1

Assume agents use both Dollars **and** Bitcoins to buy goods at t and $t + 1$. Then

$$\mathbb{E}_t \left[u'(c_{t+1}) \frac{P_t}{P_{t+1}} \right] = \mathbb{E}_t \left[\left(u'(c_{t+1}) \frac{P_t}{P_{t+1}} \right) \frac{Q_{t+1}}{Q_t} \right] \quad (1)$$

If production (consumption) is constant at $t + 1$ or if agents are risk-neutral, and if further Q_{t+1} and $1/\pi_{t+1}$ are conditionally uncorrelated, then **the Bitcoin price Q_t in Dollar is a martingale**,

$$Q_t = E_t[Q_{t+1}]$$

Bitcoin block rewards are financed by Dollar taxes

Consider two economies, which differ in the growth paths for the Bitcoin quantity.

- The central bank seeks to achieve the same path for prices.
- Quantity theory:

$$P_t y_t = D_t + Q_t B_t$$

- More Bitcoins B_t means less D_t , keeping everything else the same.
- Same equilibrium can obtain, otherwise.

S-FV-U, “CBDC: When Price & Bank Stability Collide”

In our model: Only HH, CB, projects. CB is financial intermediary.

Key Mechanism

- Nominal Diamond-Dybvig (1983) model for a CB and its CBDC.
- Central bank can always deliver on its nominal obligations.
- But: CB runs can happen: “spending run” on available goods.

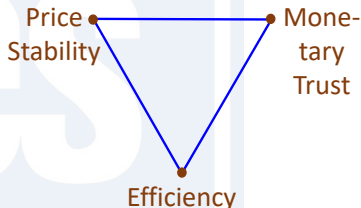
Three competing objectives:

- 1 Traditional CB objective: commitment to **Price Stability**
- 2 Social optimum, optimal risk sharing: **Efficiency**
- 3 Absence of runs, financial stability: **Monetary Trust**

Key Result:

CBDC Trilemma

Of the three objectives, the central bank can only achieve two.



The model: the real portion is Diamond-Dybvig, 1983

- time $t = 0, 1, 2$.
- Continuum $[0, 1]$ of agents:
 - ▶ $t = 0$: symmetric, endowed with one unit of a real good
 - ▶ $t = 1$: types reveal: “impatient” λ , “patient” $1 - \lambda$.
Impatient agents: have to consume in $t = 1$.
 - ▶ $u(\cdot)$ strictly increasing, concave, RRA greater than one,
 $-x \cdot u''(x)/u'(x) > 1$.
- Real Technology:
 - ▶ long term: $1 \rightarrow 1 \rightarrow R$
 - ▶ storage $t = 1 \rightarrow t = 2$, available to all: $1 \rightarrow 1$
- **Optimal solution:**

$$\max \lambda u(x_1) + (1 - \lambda)u(x_2) \quad \text{s.t.} \quad \lambda x_1 + (1 - \lambda)\frac{x_2}{R} = 1$$

Unique solution, where $u'(x_1^*) = Ru'(x_2^*)$

- With that: $x_1^* > 1$. (Diamond and Dybvig, 1983)

The model: the nominal portion introduces CBDC.

- $t = 0$:
 - Agents sell goods to CB for M CBDC units in $t = 1$.
 - CB: invests all received real goods in projects.
- $t = 1$:
 - Agents learn type. Impatient agents spend M . Patient agents may. Total fraction: $\lambda \leq n \leq 1$.
 - CB **observes agg. spending fraction n** .
 - CB liquidates fraction $y = y(n) \in [0, 1]$ of projects.
 - CB sells goods y . Market clearing price P_1 .
- $t = 2$:
 - Remaining agents spend $(1 + i(n))M$.
 - CB sells remaining project payoffs $R(1 - y)$
 - Market clearing price P_2 .

Definition 1

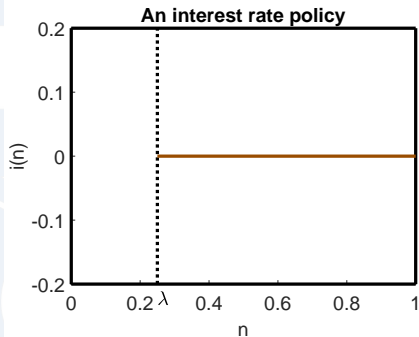
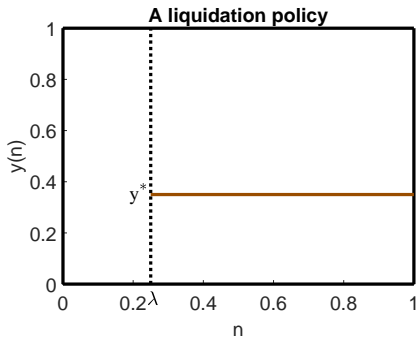
A **central bank policy** is a triple $(M, y(\cdot), i(\cdot))$, where $y : [0, 1] \rightarrow [0, 1]$ is the central bank's liquidation policy for every observed fraction n of spending agents, and $i : [0, 1] \rightarrow [-1, \infty)$ is the nominal interest rate policy.

A (boring) example for a central bank policy

A **central bank policy** is a triple $(M, y(\cdot), i(\cdot))$:

$y(n)$

$i(n)$



Set M so that $P_1 = 1$ clears the market, if $n = \lambda$ agents spend in $t = 1$.

Market Clearing

$$\begin{aligned}nM &= P_1 y(n) \\ (1-n)(1+i(n))M &= P_2 R(1-y(n)),\end{aligned}$$

$\Rightarrow n, y(n), i(n)$ pin down the price levels P_1, P_2 .

$$P_1(n) = \frac{nM}{y(n)} \quad \text{and} \quad P_2(n) = \frac{(1-n)(1+i(n))M}{R(1-y(n))}$$

Note: $P_2(n)$ can be “anything” per $i(n)$, but $i(n)$ does not affect $P_1(n)$.

Real allocation: only depends on n via $y(n)$:

$$x_1(n) = \frac{M}{P_1} = \frac{y(n)}{n} \quad \text{and} \quad x_2(n) = \frac{(1+i(n))M}{P_2} = \frac{1-y(n)}{1-n} R$$

Given n , patient agents run iff $x_1(n) \geq x_2(n)$.

Objective 2: Optimal Risk Sharing

The social optimum (x_1^*, x_2^*) is an equilibrium, if $y(\lambda) = y^* = \lambda x_1^*$.

Objective 3: Absence of Runs

A Run on the Central Bank is a Spending Run:

Definition 2

A **run** occurs if $n > \lambda$: patient agents also spend.

CBDC loses its 'store of value' function.

- Patient agents purchase goods instantaneously even though they do not need to consume
- Enable future consumption by storing toilet paper and other goods at home rather than storing value in form of CBDC
- Trust in monetary system and CBDC evaporates.
- Monetary instability.
- Compare to:
 - ▶ temporary pandemic stockouts.
 - ▶ hyperinflations.
 - ▶ currency crises.

No run

t=0



t=1



t=2



Run

t=0



t=1

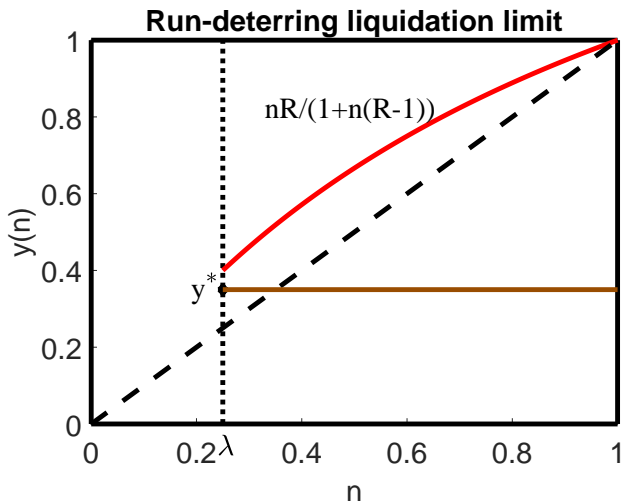


t=2



Run-Proof Policies

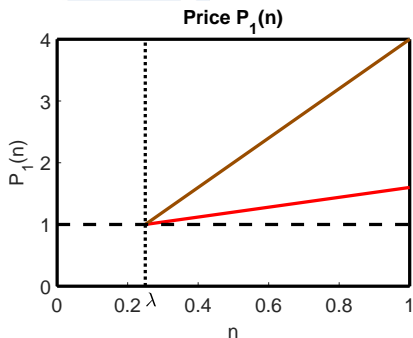
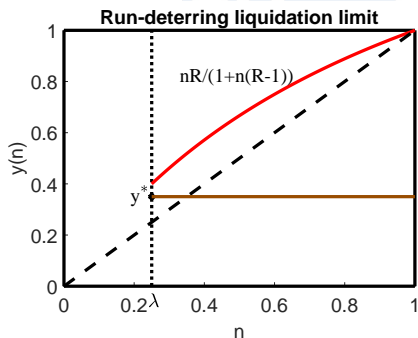
The policy is “run-proof”, if $n \neq \lambda$ is “off equilibrium”, i.e. if $x_1(n) < x_2(n)$ for all n , i.e. $y(n) < \bar{y}(n) = nR/(1 + n(R - 1))$.



Example: the policy $y(n) \equiv y^*$ is run-proof.

Run-Proof Policies and Price Implications

E.g. for $y(n) \equiv y^*$, we have $P_1(n) = n \frac{M}{y^*}$



- These two policies violate the **price stability** objective for $P_1(n)$.
- The problem only arises “off equilibrium.”
- Commitment-issue / credibility / sub-game perfection: should $n \neq \lambda$ arise, a price-stability oriented Central Bank may not stick to the “threat” of letting the price P_1 move far from the target.
- Remark: objective for $P_2(n)$ can always be achieved via $i(n)$.

Objective 1: Price Stability

Definition 3

- 1 A central bank policy is **fully price stable**, if it achieves $P_1(n) \equiv \bar{P}$ for all n .
- 2 A central bank policy is **partially price stable**, if it achieves **either** $P_1(n) = \bar{P}$ **or** there is full liquidation, $y(n) = 1$, for all n .

(In the paper: extend to period 2, pick the right interest rate policy.)

Recall Market Clearing:

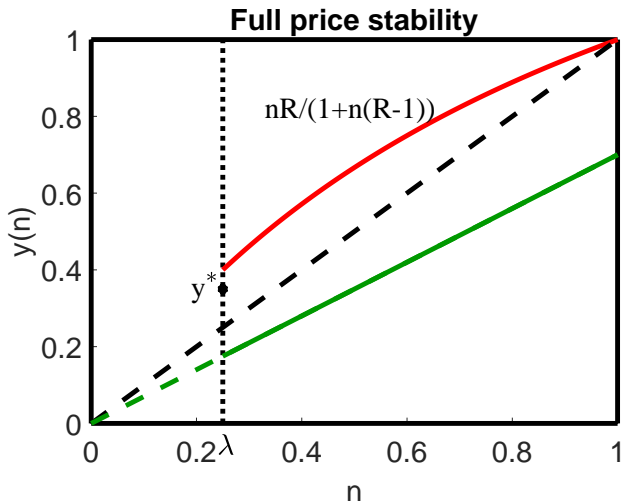
$$P_1(n) = \frac{nM}{y(n)}$$

Therefore,

fully price stable: $y(n) = \frac{nM}{\bar{P}}$

partially price stable: $y(n) = \min\left\{\frac{nM}{\bar{P}}\right\}$

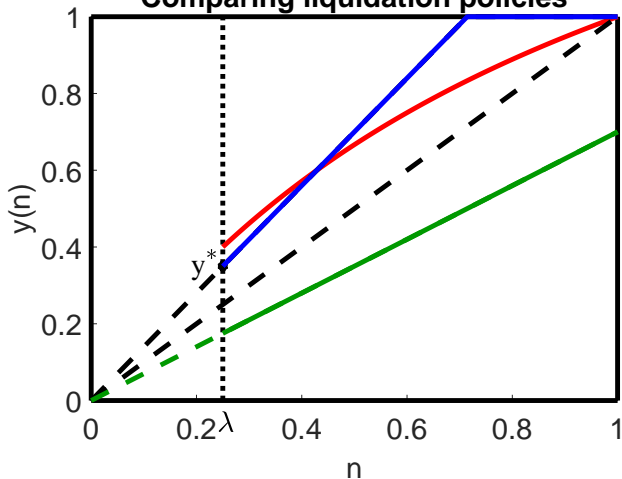
Full price stability



- Prices are fully stable and runs are avoided on green line.
- **But:** no longer efficient at $n = \lambda$!
- At best: green line = 45 degree line.

Partial price stability

Comparing liquidation policies



- Prices partially stable and efficiency on blue line.
- **But:** no longer run-proof. Runs may happen!

Conclusions

Two papers

- 1 Theory of Cryptocurrency Pricing: **Schilling-Uhlig, “Some Simple Bitcoin Economics”, JME 2019.**
- 2 Central bank digital currencies: **Schilling - Fernández-Villaverde - Uhlig, “CBDC: when Price and Bank Stability Collide”.**

... but a lot more can and should be done! Wide-open canvas.

THANKS!



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